

# CPE 390: Microprocessor Systems

Spring 2018

## Lecture 5

# Assembly Programming: Arithmetic

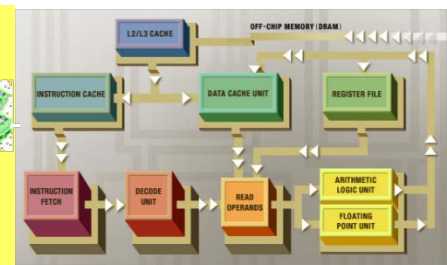
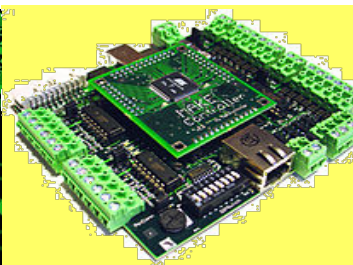
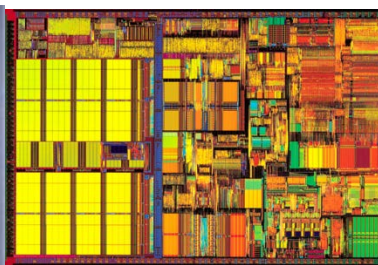
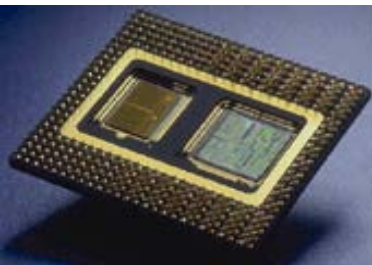
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Adapted from HCS12/9S12 An Introduction to Software and Hardware Interfacing Han-Way Huang, 2010

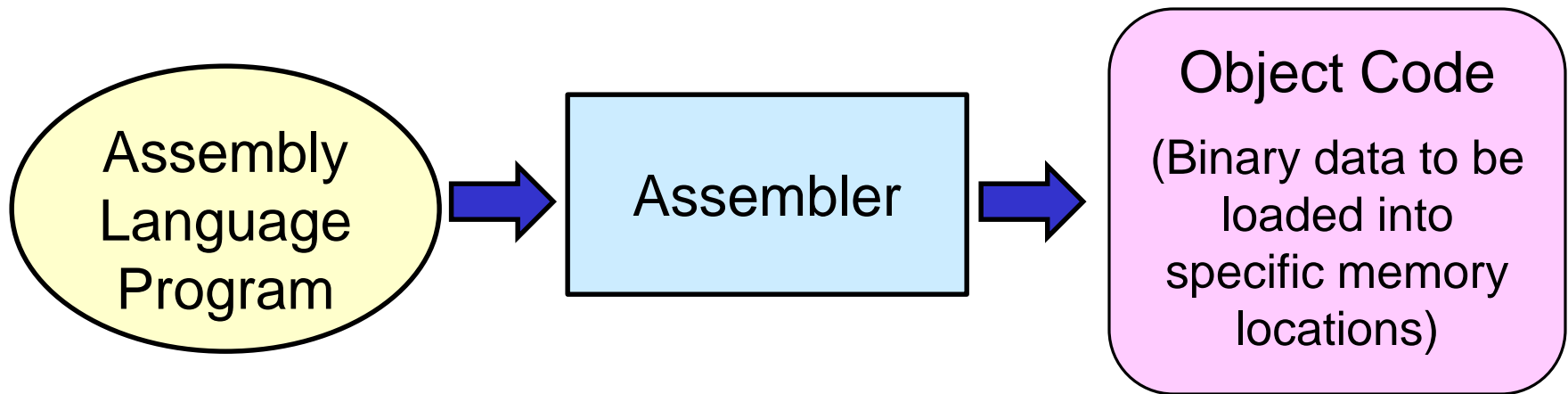


# Try These...

1. What is  $27_{10}$  in 8-bit binary?
2. What is  $-27_{10}$  in 8-bit binary?
3. What is %10011010 (unsigned) in decimal?
4. What is %10011010 (signed) in decimal?
5. What is %10011010 in hex?
6. What is %10101101 + %00100111 in binary (unsigned)
7. What is %10101101 + %00100111 in binary (signed)
8. What is  $299_{10}$  in 16-bit hex?
9. What is \$1A3F in decimal?
10. What is \$39C2 + \$A175 in hex?

# What is Assembly Language?

- Assembly Language (assembly code) allows a programmer to specify machine code *instructions* and *data* that should be loaded into microprocessor memory prior to program execution.
  - Machine code *instructions* are specified using mnemonics and address labels
  - *Data* represents initial values of program variables



- Assembler translates assembly code mnemonics & symbols into raw binary data to be loaded into microprocessor memory

# Structure of a HCS12 Assembly Program

*You will find three kinds of statements in assembly program:*

- **Assembler Directives**
  - Define data and symbols
  - Reserve and initialize memory locations
  - Set assembler and linking conditions
  - Specify output format
  - Specify end of program
- **Assembly Language Instructions**
  - mnemonic representation of HCS12 machine code instructions
- **Comments**
  - Explanation and documentation

# Program Structure: Example Code

*label*

results:

inc\_value:

din:

ORG \$800

DS.B 4

EQU \$30

DC.B \$1122

*assembler directive*

;reserve 4 bytes for result

;symbol to represent data

;label & initialize data

*comment*

*assembly  
language  
instruction*

ORG \$900

ldd din

subd #10

std results

adda #inc\_value

std results+2

END

;program begins

;load \$1122 into D

;program ends

*opcode*

std

results

*operand*

# Fields of an HCS12 Instruction

loop:    **adda**    # $\$40$     ;add  $\$40$  to accumulator A

- **Label Field**
  - optional: usually starts from first column
  - start with a letter followed by letters, digits or ( \_ or .)
  - can start any column if ended with a colon :
- **Operation (Opcode) Field**
  - mnemonic machine code instructions or assembler directive
  - is separated from label or beginning of line by at least one space
- **Operand Field**
  - operands for instructions or arguments for assembler directives
  - separated from operation field by at least one space
- **Comment Field**
  - optional: starts with ;
  - separated from operation/operand field by at least one space
  - a line that starts with \* or ; is a comment

# Some Assembler Directives

- **END**
  - Ends program to be processed by assembler
  - Any statement after **END** is ignored
- **ORG**
  - Assembler uses a location counter to keep track of current memory location
    - where next machine code byte or data byte should be placed
  - ORG directive sets a new value into the location counter
  - for example:

ORG	\$1000	→	<i>opcode</i>	\$C6	\$1000
ldab	#\$FF		<i>operand</i>	\$FF	\$1001

will place the opcode byte for the “ldab” instruction at memory address \$1000

# Initialize Memory Directives

- **DC.B** (define constant byte)
  - define value of byte (or bytes) at current memory location
  - location counter is updated to point to next byte address
  - value can be specified by expression
- **DC.W** (define constant word)
  - define value of 2-byte word(s)

array:    ORG     \$4000  
          DC.B    \$A3  
          DC.B    \$11, \$22, \$33  
          DC.W    \$1234, array-2



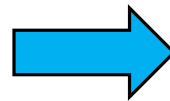
\$A3	\$4000
\$11	\$4001
\$22	\$4002
\$33	\$4003
\$12	\$4004
\$34	\$4005
\$3F	\$4006
\$FE	\$4007



# Initialize Text String

- **DC.B** can also be used to define and load a string of ascii characters
  - string specified using quotes (“”)
  - each character represented by one-byte ascii code

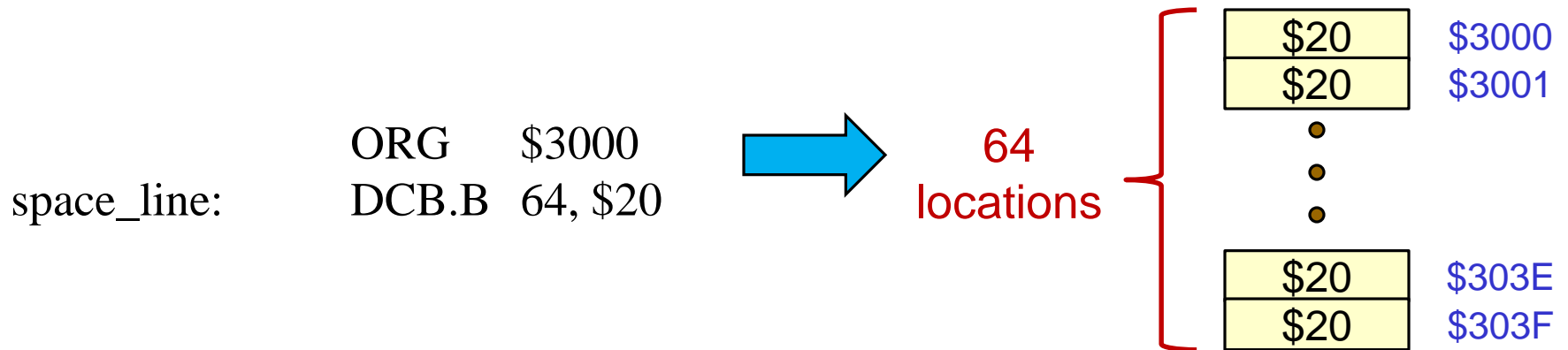
```
msg:   ORG    $1000
       DC.B   "enter y/n:"
```



\$1000	\$65	e
\$1001	\$6E	n
\$1002	\$74	t
\$1003	\$65	e
\$1004	\$72	r
\$1005	\$20	<sp>
\$1006	\$79	y
\$1007	\$2F	/
\$1008	\$6E	n
\$1009	\$3A	:

# Initialize Memory Directives

- **DCB.B** (define constant block of bytes)
  - fill a block of memory locations with same one-byte value
  - syntax is: **DCB.B count ,value**
  - **value** is optional – default value is \$00
- **DCB.W** (define constant block of words)
  - fill a block of memory with same two-byte value



# Reserve Memory Directives

- **DS.B** (define storage byte)
  - reserves (and optionally labels) number of bytes at current memory location
  - location counter is updated to point to next byte address following the reserved space
  - reserved locations are not initialized
- **DS.W** (define storage word)
  - reserves, and optionally labels: (# words X 2) bytes

```
                ORG    $1400
buffer: DS.B    $100    ;reserves 256 bytes of memory from $1400
                        ;to $14FF with "buffer" labeling first byte
wbuf:  DS.W    20      ;reserves 20 words (40 bytes) of memory from
                        ;$1500 to $1527 with "wbuf" labeling first byte
```

# Equate Directive

- **EQU** (equate)
  - assigns a value (rather than a memory address) to a label
  - does not affect memory contents

```
loop_count:    EQU    16
               ...
               ldaa   #loop_count    ;load 16 into accumulator A
```

# Assembler Directive Examples

- Show the contents of memory resulting from the following assembler directives:

```
                ORG    $4800
xyz:            EQU    24
abc:            DC.B   $20, 16
                DC.W   $21, $1ACD

res:            DS.B   3
                DC.B   "bcd"
                DC.W   abc+xyz
```

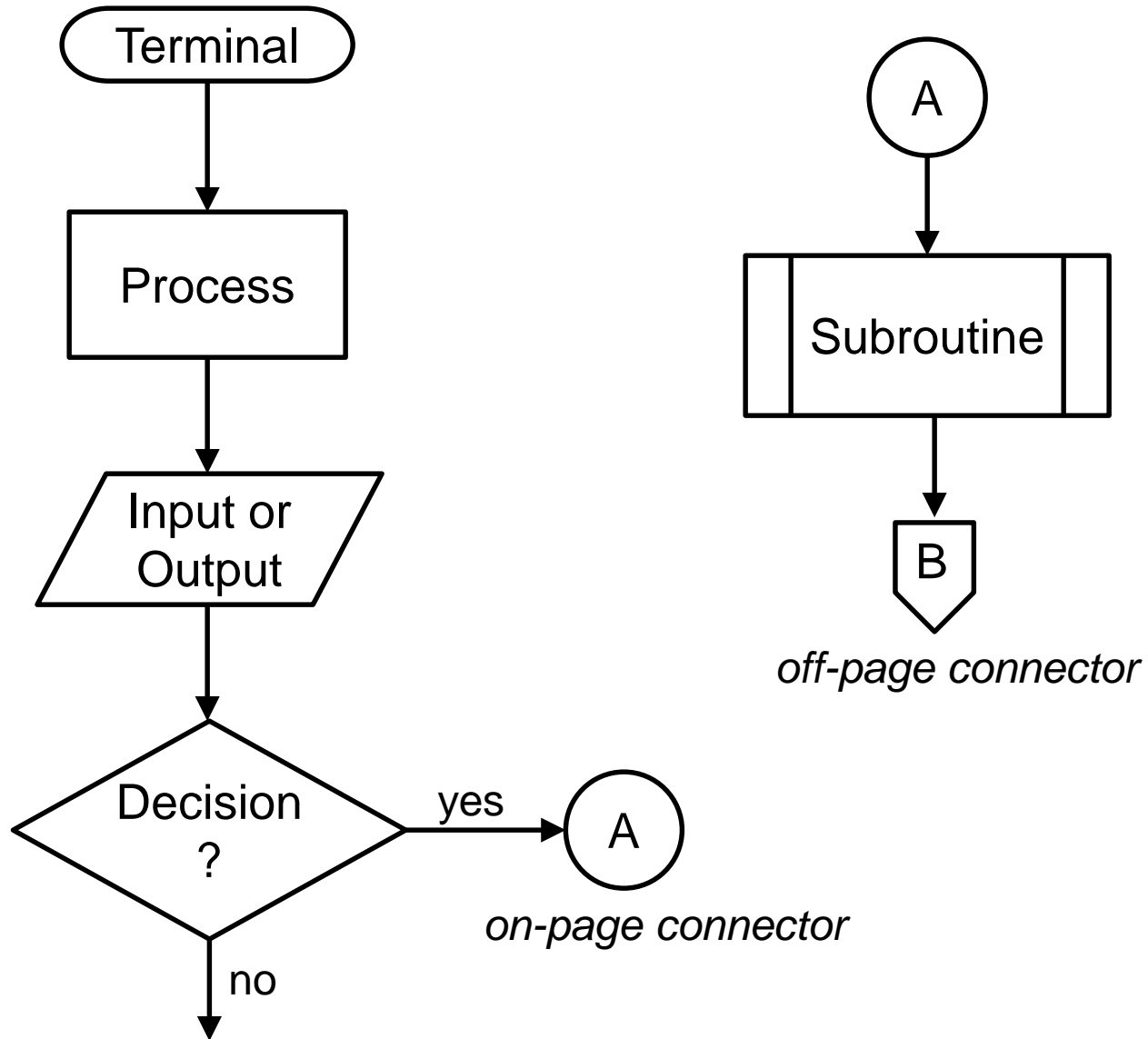
- Show results as a table:

label	address	data

# Software Development Process

- Problem definition:
  - Identify precisely what needs to be done
- Develop a plan or algorithm
  - computational procedure that takes a set of inputs and produces required outputs
  - may be expressed as a set of steps that need to be performed
  - may include iteration and sub-procedures or subroutines
  - need to specify data structures that may be required
  - algorithm may be expressed in pseudo-code (e.g.  $A \leftarrow A+1$ )
  - algorithm code may be expressed in [flow-chart](#)
- Programming
  - convert computational (or flowchart) steps into executable statements and data structures in target language
- Program testing & debugging
- Program maintenance

# Flow-Chart Symbols



# Programs to do Simple Arithmetic

- Write a program starting at memory location \$1500 to subtract the contents of memory location \$1005 from the sum of memory locations \$1000 and \$1002 and store the difference at \$1010.

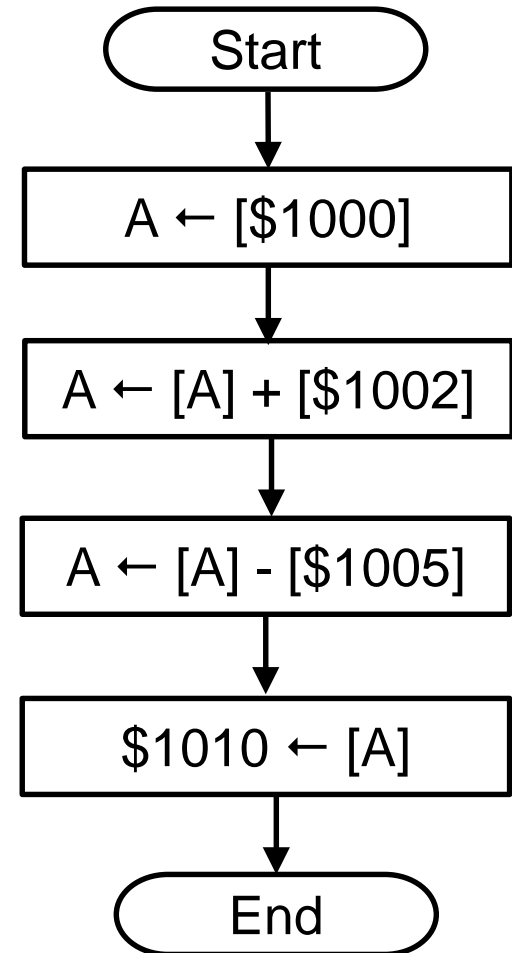
- Solution:

**Step1:** Load contents of memory loc. \$1000 into A

**Step2:** Add contents of memory loc. \$1002 to A

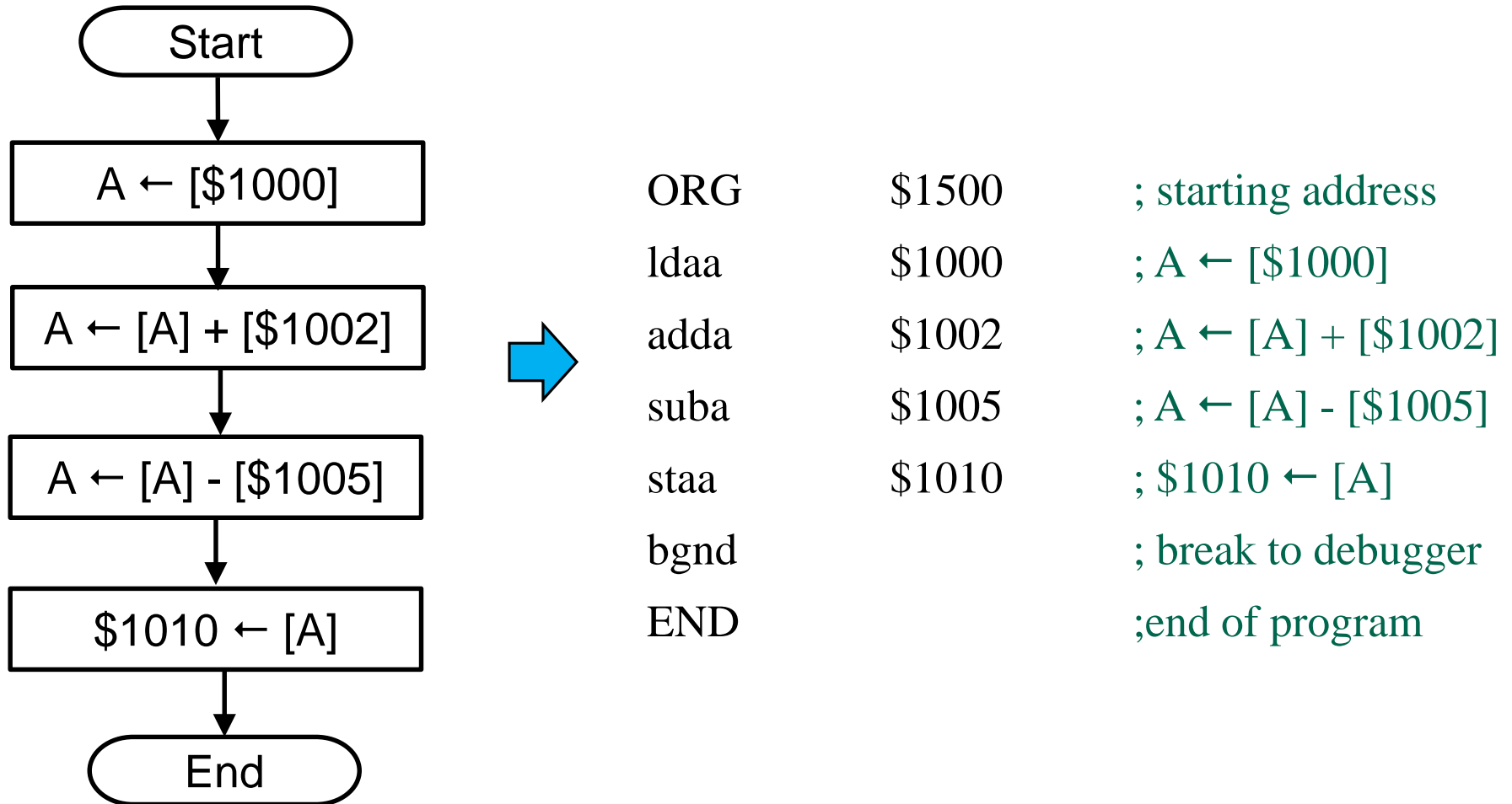
**Step3:** Subtract contents of memory loc. \$1005 from A

**Step4:** Store contents of A to memory loc. \$1010





# Algorithm to Assembly Code



# More on Arithmetic (Add/Sub)

- We know how to add 8-bit quantities using A or B accumulator:

```
ldaa    $1000    ; add 8-bit data in $1000
adda    $1001    ; to 8-bit data in $1001
staa    $2000    ; and store 8-bit result in $2000
```

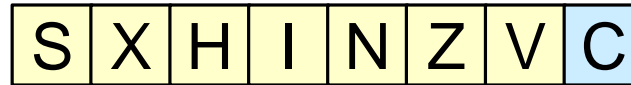
- and we can add 16-bit quantities using D accumulator:

```
ldd     $1000    ; add 16-bit data in $1000:$1001
addd    $1002    ; to 16-bit data in $1002:$1003
std     $2000    ; and store 16-bit result in $2000:$2001
```

- How can we add quantities of greater precision (e.g. 24-bit) ?

# Carry/Borrow Flag

- Carry is LSB of CCR



- Carry acts like the  $n^{\text{th}}$  result bit when doing  $n$ -bit add/sub
- Carry set to '1' whenever addition generates carry-out
- Carry set to '1' whenever subtraction requires borrow-out

ldaa #20
adda #30
20
<u>+30</u>
0 50

ldaa #93
adda #8B
93
<u>+8B</u>
1 1E

ldaa #96
suba #3A
96
<u>-3A</u>
0 5C

ldaa #27
suba #54
27
<u>-54</u>
1 D3

- also 16-bit:

ldd #A200
add #7000
A200
<u>+7000</u>
1 1200

# Multi-precision Addition

- Carry bit allows us to do multi-precision arithmetic
  - i.e. arithmetic on numbers whose precision is greater than that of the ALU
- For example how do we add \$59A183 to \$5482DB ?
- Solution:
  - Step1:** Add \$A183 to \$82DB and remember carry
  - Step2:** Store 16-bit result as least significant two bytes of answer
  - Step3:** Add \$59 to \$54 with carry from step 1
  - Step4:** Store 8-bit result as MSbyte of answer

# 32-bit Arithmetic

- Write a program at memory location \$4000 to add 4-byte numbers that are stored at \$1000~\$1003 and \$1004~\$1007 and store the sum at locations \$1010~1013

# Multiplication

Mnemonic	Function	Operation
emul	Unsigned 16 x 16 multiply	$Y:D \leftarrow [D] \times [Y]$
emuls	Signed 16 x 16 multiply	$Y:D \leftarrow [D] \times [Y]$
mul	Unsigned 8 x 8 multiply	$D \leftarrow [A] \times [B]$

- Write an instruction sequence to multiply (unsigned) register X by register Y and store result in \$1000~\$1003

```
tfr      x,d      ;transfer X operand into D
emul                    ;perform multiplication
sty      $1000    ;save upper 16-bits of product
std      $1002    ;save lower 16-bits of product
```

# Multiplication

Mnemonic	Function	Operation
emul	Unsigned 16 x 16 multiply	$Y:D \leftarrow [D] \times [Y]$
emuls	Signed 16 x 16 multiply	$Y:D \leftarrow [D] \times [Y]$
mul	Unsigned 8 x 8 multiply	$D \leftarrow [A] \times [B]$

- n1 and n2 are signed 16-bit integers and n3 is a 32-bit signed integer. Use assembler directives to reserve space for n1, n2 and n3 at memory locations \$1000, \$1002 and \$1004 respectively.

Write code starting at \$4000 to perform:  $n3 = n1 * n2$

# Division

Mnemonic	Function	Operation
ediv	Unsigned 32 by 16 divide	$Y \leftarrow [Y]: [D] \div [X]$ $D \leftarrow remainder$
edivs	Signed 32 by 16 divide	$Y \leftarrow [Y]: [D] \div [X]$ $D \leftarrow remainder$
idiv	Unsigned 16 by 16 divide	$X \leftarrow [D] \div [X]$ $D \leftarrow remainder$
idivs	Signed 16 by 16 divide	$X \leftarrow [D] \div [X]$ $D \leftarrow remainder$

- Remember that if divisor > dividend, then quotient = 0



# Integers Math and Precision

- When performing integer arithmetic (especially multiplication and division), important to keep track of potential size of results to avoid overflow and/or loss of precision
- Suppose we want to calculate: 
$$\frac{1200 \times 2500}{1150}$$
- Does the order of the operations matter?
- Correct answer is **2608.6956**
  - (but we can only do integer arithmetic)
- If I do multiply first (emul followed by ediv), I get 2608
- If I do divide first ( ediv followed by emul), I get 2500

# Rounding

- If we take the quotient as being the answer to a divide operation (and ignore the remainder), the result is truncated to the closest integer that is less than the correct answer ( 2608 instead of 2608.6956)
- A better result would be to round to the nearest integer (2609). This can be achieved by adding half of the divisor to the dividend before executing the divide operation:

$$\textit{rounded quotient} = \frac{\textit{dividend} + (\textit{divisor}/2)}{\textit{divisor}}$$

- This effectively adds 0.5 to the answer before truncation.

# Integer Precision: Example

- Multiply the unsigned 16-bit number in locations \$1000:\$1001 by 1.414 (approx. to  $\sqrt{2}$ ), truncating result to nearest integer.
  
- Multiply the unsigned 16-bit number in locations \$1000:\$1001 by 1.414 (approx. to  $\sqrt{2}$ ), rounding result to nearest integer.