# Exact Reconstruction Analysis of Log-Sum Minimization for Compressed Sensing

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Abstract—The fact that fewer measurements are needed by log-sum minimization for sparse signal recovery than the  $\ell_1$ -minimization has been observed by extensive experiments. Nevertheless, such a benefit brought by the use of the log-sum penalty function has not been rigorously proved. This paper provides a theoretical justification for adopting the log-sum as an alternative sparsity-encouraging function. We prove that minimizing the log-sum penalty function subject to Az = y is able to yield the exact solution, provided that a certain condition is satisfied. Specifically, our analysis suggests that, for a properly chosen regularization parameter, exact reconstruction can be attained when the restricted isometry constant  $\delta_{3K}$  is smaller than one, which presents a less restrictive isometry condition than that required by the conventional  $\ell_1$ -type methods.

*Index Terms*—Compressed sensing, iterative reweighted algorithms, log-sum minimization.

#### I. INTRODUCTION

T HE problem of compressed sensing involves the recovery of a high dimensional sparse signal from a small number of measurements [1], [2]. The canonical form of this problem can be presented as

$$\min_{\boldsymbol{x}} \|\boldsymbol{z}\|_0 \quad \text{s.t.} \quad \boldsymbol{A}\boldsymbol{z} = \boldsymbol{A}\boldsymbol{x} \tag{1}$$

where y = Ax denotes the acquired measurements,  $A \in \mathbb{R}^{m \times n}$  is the sampling matrix with  $m \ll n$ , and  $||z||_0$  stands for the number of the nonzero components of z. It is well-known that any K-sparse vector x can be exactly recovered via (1) if  $\delta_{2K} < 1$ , where  $\delta_{2K}$  is the restricted isometry constant associated with the measurement matrix A, which is defined as the smallest constant such that

$$(1 - \delta_{2K}) \|\boldsymbol{z}\|_{2}^{2} \le \|\boldsymbol{A}\boldsymbol{z}\|_{2}^{2} \le (1 + \delta_{2K}) \|\boldsymbol{z}\|_{2}^{2}$$
(2)

holds for all 2K-sparse vectors [1]. The optimization (1), however, is a non-convex and NP-hard problem that has computational complexity growing exponentially with the signal dimension n. Thus, alternative sparsity-promoting functionals which

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are more computationally efficient in finding the sparse solution are desirable. The most popular alternative is to replace the  $\ell_0$ -norm in (1) with  $\ell_1$ -norm, which leads to a convex optimization problem that can be solved efficiently. Over the past decade, the use of the  $\ell_1$ -norm as a sparsity-promoting functional for sparse signal recovery has been extensively studied [1]–[6]. It has been shown that  $\ell_1$  minimization allows recovery of sparse signals from only a few measurements. Nevertheless, as compared with (1),  $\ell_1$ -type methods generally require a more restrictive condition for exact signal reconstruction. It is therefore natural to seek an alternative which can bridge the gap between the  $\ell_0$  and  $\ell_1$  minimization. One such alternative is the log-sum penalty function. Replacing the  $\ell_0$ -norm with this sparsity-encouraging functional leads to:

$$\min_{\boldsymbol{z}} \quad L(\boldsymbol{z}) = \sum_{i=1}^{n} \log(|z_i| + \epsilon) \quad \text{s.t.} \quad \boldsymbol{A}\boldsymbol{z} = \boldsymbol{A}\boldsymbol{x} \quad (3)$$

where  $\epsilon > 0$  is a positive parameter to ensure that the function is well-defined. Such a log-sum penalty function was originally introduced in [7] for basis selection and has gained increasing attention recently. It was shown in [8], [9] that by resorting to the bound optimization technique, minimizing (3) can be formulated as an iterative reweighted  $\ell_1$ -minimization process which iteratively minimizes a reweighted  $\ell_1$  function. In a series of experiments [8], the iterative reweighted algorithm presents uniform superiority over the conventional  $\ell_1$ -type methods in the sense that substantially fewer measurements are needed for exact recovery. In fact, is was shown in [10] that when  $\epsilon =$ 0, the log-sum penalty function is essentially the same as the  $\ell_0$ -norm. Hence it can be expected that the above regularized log-sum penalty function behaves like the  $\ell_0$ -norm when  $\epsilon$  is small. Nevertheless, such a benefit brought by the use of the regularized log-sum penalty function in sparse signal recovery has not been rigorously proved so far. In the following, we conduct an in-depth investigation of the optimization (3). Our study will provide a rigorous justification for (3) and the iterative reweighted method.

In addition to the log-sum minimization being considered in this paper, another effective sparsity-promoting strategy is to replace the  $\ell_0$ -norm with the  $\ell_p$ -norm (0 ). The $oretical analyses conducted in [11]–[13] prove that <math>\ell_p$ -minimization enjoys a nice theoretical guarantee:  $\ell_p$ -minimization requires fewer measurements than  $\ell_1$ -type methods for sparse signal exact reconstruction. We show that a similar theoretical guarantee is available for the minimization of log-sum penalty function as well.

#### **II. CONDITION FOR EXACT RECONSTRUCTION**

We provide theoretical analysis concerning a sufficient condition under which the solution to (3) equals to the true signal *x*. Our analysis reveals a relation between the restricted isometry constant  $\delta_{3K}$  and the regularization parameter  $\epsilon$ . The main result is summarized as follows.

Theorem 1: Let A be an  $m \times n$  matrix, and  $x \in \mathbb{R}^n$  be a sparse vector with K non-zero entries. Define

$$x_{\min} \triangleq \min_{i \in T_0} |x_i| \qquad \beta \triangleq \frac{\sqrt{1 - \delta_{3K}}}{2\sqrt{K}(L - 1)} \tag{4}$$

where  $T_0$  is the support of  $\boldsymbol{x}, L \triangleq \lceil \frac{n-K}{2K} \rceil, \lceil u \rceil$  denotes the ceiling operator that gives the smallest integer no smaller than u, and  $\delta_{3K}$  is the restricted isometry constant of the measurement matrix  $\boldsymbol{A}$ . If we have

$$\frac{x_{\min}}{\epsilon} > \min_{\{\tau: 0 < \tau < 1\}} f(\tau) \tag{5}$$

where

$$f(\tau) \triangleq \max\left\{\frac{(1/\beta) - 1}{1 - \tau}, \frac{2}{(\beta\tau)^2}\right\}$$
(6)

then the global minimizer of (3) is exactly equal to  $\boldsymbol{x}$ .

*Remark 1:* The above condition (5) involves a search of  $\tau$  over the region (0, 1). Notice that the first term inside the max (one out of two) operator in (6) tends to infinity when  $\tau \to 1$ , while the second term becomes arbitrarily large when  $\tau \to 0$ . Hence the value of  $\tau$  minimizing  $f(\tau)$  should lie somewhere in between. Nevertheless, when we are trying to select an appropriate parameter  $\epsilon$  to satisfy (5), an arbitrary value of  $\tau \in (0, 1)$  can be considered. As long as  $x_{\min}/\epsilon > f(\tau)$  is satisfied for a particular choice of  $\tau_0 \in (0, 1)$ , we can guarantee that the condition (5) is automatically satisfied since

$$f(\tau_0) \ge \min_{\{\tau: 0 < \tau < 1\}} f(\tau) \qquad \tau_0 \in (0, 1)$$
(7)

For example, if we select  $\tau = 10/11$ , then we only need to choose an  $\epsilon$  to ensure

$$\frac{x_{\min}}{\epsilon} > \max\left\{11\left(\frac{1}{\beta} - 1\right), 2\left(\frac{1.1}{\beta}\right)^2\right\}$$
(8)

We would like to emphasize that the condition (5) is a sufficient condition for exact reconstruction. When  $\epsilon$  becomes arbitrarily large, the condition (5) will not be satisfied. Nevertheless, in this case, exact reconstruction is still possible. Note that when implementing the iterative reweighted  $\ell_1$ -minimization algorithm, all weights are roughly identical. Hence the log-sum penalty function behaves like  $\ell_1$ -minimization for an arbitrarily large  $\epsilon$ .

*Remark 2:* Theorem 1 provides a sufficient condition which guarantees an exact reconstruction via solving (3). A close examination of the condition (5) reveals that the regularization parameter  $\epsilon$  has an inverse relationship with the restricted isometry constant  $\delta_{3K}$ , that is, a larger  $\delta_{3K}$  results in a smaller  $\epsilon$  and vice versa. In particular, when  $\delta_{3K} \rightarrow 1$ , accordingly we should have  $\epsilon \rightarrow 0$  in order to ensure the condition (5) is met. One the other hand, for any  $0 < \delta_{3K} < 1$ , we can always find a sufficiently small  $\epsilon$  such that the condition (5) is satisfied. Hence we can ensure that, when  $\delta_{3K} < 1$ , any *K*-sparse signal can be exactly recovered via (3) for a properly chosen  $\epsilon$ . Recalling that for  $\ell_1$ -minimization methods, the condition for exact reconstruction is given by  $\delta_{3k} + 3\delta_{4k} < 2$  (see [3]). Since we have  $\delta_{3k} < \delta_{4k}$ , the condition  $\delta_{3k} + 3\delta_{4k} < 2$  implies  $\delta_{3k} < 0.5$ . We see that our

result presents a less restrictive isometry condition than that of the  $\ell_1$ -minimization methods. This also explains why the use of the log-sum penalty function turns out to be a better alternative to the  $\ell_0$  than the  $\ell_1$ -norm.

*Remark 3*: We note that (3) is a non-convex optimization problem and there is no guarantee that the iterative reweighted algorithm will converge to a global minimum of (3). Nevertheless, we can improve the probability of finding a global minimizer by starting from a number of different initialization points and choosing the converged point that achieves the minimum objective function value. Also, empirical studies suggest that the iterative reweighted algorithm is more likely to converge to an undesirable local minimum when  $\epsilon \rightarrow 0$ . To address this issue, similarly to [14], we can use a monotonically decreasing sequence  $\{\epsilon^{(t)}\}\$  in updating the weighting parameters. For example, at the very first beginning,  $\epsilon^{(t)}$  can be set to a relatively large value, say 1, in order to provide a stable coefficient estimate. We then gradually reduce the value of  $\epsilon^{(t)}$  in the subsequent iterations until  $\epsilon^{(t_k)}$  attains a value such that (5) is met. Such a process actually amounts to solving the following optimization

$$\min_{\boldsymbol{z}} \quad L(\boldsymbol{z}) = \sum_{i=1}^{n} \log(|z_i| + \epsilon^{(t_k)}) \quad \text{s.t.} \quad \boldsymbol{A}\boldsymbol{z} = \boldsymbol{A}\boldsymbol{x} \quad (9)$$

Previous iterations that use  $\epsilon^{(t)} \neq \epsilon^{(t_k)}$  can be considered a procedure looking for a good initialization point. Numerical results demonstrate that this approach significantly improves the ability of avoiding undesirable local minima.

## III. PROOF OF THEOREM 1

Suppose that  $x^*$  is the solution of (3). Let  $h \triangleq x^* - x$ . Clearly, the residual vector h lies in the null space of A, i.e.

$$\boldsymbol{A}\boldsymbol{h} = \boldsymbol{0} \tag{10}$$

Meanwhile, since  $x^*$  is the global minimizer of (3), we have

$$L(\boldsymbol{x}^*) = L(\boldsymbol{x} + \boldsymbol{h}) \le L(\boldsymbol{x}) \tag{11}$$

We wish to prove that, given (5), there does not exist a nonzero vector h which satisfies conditions (10)–(11) simultaneously. Otherwise the global minimizer of (3) is unequal to the original sparse signal, i.e.  $x^* \neq x$ .

Let us first examine (10). Write  $\mathbf{h} \triangleq [h_1 \quad h_2 \quad \dots \quad h_n]^T$ . We decompose the residual vector  $\mathbf{h}$  into a sum of a set of vectors  $\{\mathbf{h}_{T_l}\}_{l=0}^L$ , where  $\mathbf{h}_{T_l}$  is a vector with its *i*th entry equal to  $h_i$  for  $i \in T_l$  and zero otherwise. The index set  $T_0$  is the support of  $\mathbf{x}$ . Also, we use  $T_0^c$  to denote the complement of the index set  $T_1$ . Without loss of generality, we assume that the index set  $T_1$  contains the indices associated with the 2K largest (in magnitude) coefficients of  $\mathbf{h}_{T_0^c}$ ,  $T_2$  corresponds to the indices of the next 2K largest (in magnitude) coefficients of  $\mathbf{h}_{T_0^c}$ , and so on. Obviously  $\{\mathbf{h}_{T_l}\}_{l=0}^L$  are 2K-sparse vectors (possibly except  $\mathbf{h}_{T_L}$ ). Our analysis will reveal the relation between  $\mathbf{h}_{T_0}$  and the largest 2K entries of  $\mathbf{h}_{T_0^c}$ . The results are summarized as follows.

*Lemma 1:* Suppose h is a vector in the null space of A, then the 2K largest (in magnitude) coefficients in  $h_{T_0^c}$  are lower bounded by

$$|h_i| \ge \beta \|\boldsymbol{h}_{T_0}\|_2 \qquad \forall i \in T_1 \tag{12}$$

Proof: See Appendix A

The above result suggests that for any nonzero vector  $\boldsymbol{h}$  in the null space of  $\boldsymbol{A}$ , the largest 2K entries in  $\boldsymbol{h}_{T_0^c}$  cannot be made arbitrarily small relative to  $\boldsymbol{h}_{T_0}$ .

Now let us consider the second condition (11). Decomposing the indices into two sets, the condition can be re-expressed as

$$\sum_{i \in T_0} \log\left(\frac{|x_i + h_i| + \epsilon}{|x_i| + \epsilon}\right) \le \sum_{i \in T_0^c} \log\left(\frac{\epsilon}{|h_i| + \epsilon}\right)$$
(13)

Based on (12) and (13), we now derive a new condition h has to satisfy. Define

$$\alpha_i \triangleq |x_i|/\epsilon \quad \forall i \in T_0 \qquad \alpha_{\min} \triangleq \min_{i \in T_0} \alpha_i$$
$$\eta_i \triangleq |h_i|/\epsilon \quad \forall i \in T_0 \qquad \eta_{\max} \triangleq \max_{i \in T_0} \eta_i$$

The term on the left-hand side (LHS) of (13) can be lower bounded by

$$\sum_{i \in T_0} \log\left(\frac{|x_i + h_i| + \epsilon}{|x_i| + \epsilon}\right) \ge \sum_{i \in T_0} \log\left(\frac{|\alpha_i - \eta_i| + 1}{\alpha_i + 1}\right)$$
(14)

while the term on the right-hand side (RHS) of (13) can be upper bounded as

$$\sum_{i \in T_0^c} \log\left(\frac{\epsilon}{|h_i| + \epsilon}\right)^{(a)} \leq \sum_{i \in T_1} \log\left(\frac{\epsilon}{|h_i| + \epsilon}\right)$$
$$\stackrel{(b)}{\leq} 2K \log\left(\frac{1}{\beta\eta_{\max} + 1}\right) \qquad (15)$$

where (a) holds valid as all terms in the summation are non-positive, and (b) follows from Lemma 1:

$$|h_i| \ge \beta \|\boldsymbol{h}_{T_0}\|_2 \ge \beta \max_{j \in T_0} |h_j| = \beta \eta_{\max} \epsilon \quad \forall i \in T_1$$
 (16)

Combining (13)–(15), we arrive at the following inequality

$$\sum_{i \in T_0} \log\left(\frac{|\alpha_i - \eta_i| + 1}{\alpha_i + 1}\right) \le 2K \log\left(\frac{1}{\beta \eta_{\max} + 1}\right) \quad (17)$$

We now prove that given (5), the above inequality (17) holds only when  $h_{T_0} = 0$  (note that the inequality (17) only involves entries in  $h_{T_0}$ ). To this objective, it suffices to prove that the converse, i.e.

$$\prod_{i \in T_0} \left( \frac{|\alpha_i - \eta_i| + 1}{\alpha_i + 1} \right) > \left( \frac{1}{\beta \eta_{\max} + 1} \right)^{2K}$$
(18)

always holds for  $h_{T_0} \neq 0$  given (5). Define  $\gamma_i \triangleq \eta_i / \alpha_i, \forall i \in T_0$ and divide the set  $\{\gamma_i\}$  into two subsets:

$$S_1 = \{i : \gamma_i \le \tau\}$$
  $S_2 = \{i : \gamma_i > \tau\}$  (19)

where  $\tau$  is a parameter of our own choice. Here we confine  $\tau$  to be a value between 0 and 1, i.e.  $0 < \tau < 1$ , to facilitate our following analysis. Clearly, we have  $S_1 \subseteq T_0$ ,  $S_2 \subseteq T_0$  and  $S_1 \cup S_2 = T_0$ . Let  $k_0$  denote the cardinality of the set  $S_1$ , and the cardinality of  $S_2$  is  $K - k_0$ . The term on the LHS of (18)

can be decomposed into a product of two terms associated with the two subsets  $S_1$  and  $S_2$ :

$$\prod_{i \in T_0} \left( \frac{|\alpha_i - \eta_i| + 1}{\alpha_i + 1} \right)$$
$$= \prod_{i \in S_1} \left( \frac{|\alpha_i - \eta_i| + 1}{\alpha_i + 1} \right) \prod_{i \in S_2} \left( \frac{|\alpha_i - \eta_i| + 1}{\alpha_i + 1} \right) \quad (20)$$

Examine the term associated with the first subset  $S_1$ . Since  $\gamma_i \leq \tau < 1$  for  $i \in S_1$ , we have  $\eta_i < \alpha_i$ . For any  $i \in S_1$ , it can be readily verified the following inequality (a) holds

$$\frac{|\alpha_i - \eta_i| + 1}{\alpha_i + 1} = \frac{\alpha_i - \eta_i + 1}{\alpha_i + 1} \stackrel{(a)}{>} \frac{1}{\beta \eta_i + 1} \ge \frac{1}{\beta \eta_{\max} + 1} \quad (21)$$

if  $\eta_i \neq 0, \forall i \in T_0$  (that is,  $h_{T_0} \neq 0$ ) and

$$\alpha_i > \frac{(1/\beta) - 1}{1 - \gamma_i} \quad \forall i \in S_1$$
(22)

Noting that  $1 - \gamma_i \ge 1 - \tau$ , the above condition (22) is guaranteed when

$$\alpha_{\min} > \frac{(1/\beta) - 1}{1 - \tau} \tag{23}$$

From (21)–(23), we see that given the condition (23), the following inequality holds

$$\prod_{i \in S_1} \left( \frac{|\alpha_i - \eta_i| + 1}{\alpha_i + 1} \right) > \left( \frac{1}{\beta \eta_{\max} + 1} \right)^{\kappa_0}$$
(24)

Consider the term associated with the subset  $S_2$  in (20). In Appendix B, we proved that the following inequality holds

$$\prod_{i \in S_2} \left( \frac{|\alpha_i - \eta_i| + 1}{\alpha_i + 1} \right) \ge \left( \frac{1}{\beta \eta_{\max} + 1} \right)^{2K - k_0}$$
(25)

when the following condition is met

$$\alpha_{\min} > \frac{2}{\left(\beta\tau\right)^2} \tag{26}$$

If  $\alpha_{\min}$  satisfies both conditions (23) and (26), i.e.

$$\alpha_{\min} > \max\left\{\frac{(1/\beta) - 1}{1 - \tau}, \frac{2}{(\beta\tau)^2}\right\}$$
(27)

by combining (24)–(25), we obtain the inequality (18). Therefore we conclude that the inequality (17) holds only when  $h_{T_0} =$ **0** (note that the inequality becomes an equality when  $h_{T_0} =$  **0**). Substituting  $h_{T_0} =$  **0** back into (13), we can quickly reach that  $h_{T_0^c} =$  **0** as well. Also, notice that the parameter  $\tau$  in (27) can take any value in (0, 1). Therefore as long as

$$\alpha_{\min} = \frac{x_{\min}}{\epsilon} > \min_{\tau \in (0,1)} \max\left\{\frac{(1/\beta) - 1}{1 - \tau}, \frac{2}{(\beta\tau)^2}\right\}$$
(28)

we have h = 0, which implies the global minimizer of (3),  $x^*$ , is exactly equal to x. The proof is completed here.

### **IV. CONCLUSIONS**

We presented a sufficient condition which guarantees that the global minimizer of (3) yields the exact reconstruction. Analyses show that when  $\delta_{3K} < 1$ , any K-sparse signal can be exactly reconstructed (3) for a properly chosen  $\epsilon$ .

## APPENDIX I PROOF OF LEMMA 1

Define a new index set  $T_{01} = T_0 \cup T_1$ . For any vector **h** in the null space of **A**, we have

$$\begin{aligned} \boldsymbol{A}\boldsymbol{h} &= \boldsymbol{A}\boldsymbol{h}_{T_{01}} + \boldsymbol{A}\boldsymbol{h}_{T_{01}^c} = \boldsymbol{0} \\ &\Rightarrow \|\boldsymbol{A}\boldsymbol{h}_{T_{01}^c}\|_2 - \|\boldsymbol{A}\boldsymbol{h}_{T_{01}}\|_2 = 0 \end{aligned} \tag{29}$$

By decomposing  $h_{T_{01}^c}$ , we reach the following inequality

$$0 = \|\boldsymbol{A}\boldsymbol{h}_{T_{01}^{c}}\|_{2} - \|\boldsymbol{A}\boldsymbol{h}_{T_{01}}\|_{2} \leq \sum_{l=2}^{L} \|\boldsymbol{A}\boldsymbol{h}_{T_{l}}\|_{2} - \|\boldsymbol{A}\boldsymbol{h}_{T_{01}}\|_{2}$$

$$\stackrel{(a)}{\leq} \sqrt{1 + \delta_{2k}} \sum_{l=2}^{L} \|\boldsymbol{h}_{T_{l}}\|_{2} - \sqrt{1 - \delta_{3K}} \|\boldsymbol{h}_{T_{01}}\|_{2}$$

$$\stackrel{(b)}{\leq} \sqrt{2}(L-1) \|\boldsymbol{h}_{T_{2}}\|_{2} - \sqrt{1 - \delta_{3K}} \|\boldsymbol{h}_{T_{0}}\|_{2}$$

$$\stackrel{(c)}{\leq} 2\sqrt{K}(L-1) |h_{T_{2}^{1}}| - \sqrt{1 - \delta_{3K}} \|\boldsymbol{h}_{T_{0}}\|_{2}$$
(30)

where (a) is a result of the restricted isometry property of A, (b) follows from the fact  $\delta_{2K} \leq 1$  and the decreasing order of  $\{\boldsymbol{h}_{T_l}\}_{l=1}^L$ , (c) comes from  $\|\boldsymbol{h}_{T_2}\|_2 \leq \sqrt{2K}|h_{T_2^1}|$ ,  $h_{T_2^1}$  is the first (meanwhile the largest in magnitude) entry in  $\boldsymbol{h}_{T_2}$ . Rearranging (30), we get

$$|h_{T_{2}^{1}}| \geq \frac{\sqrt{1 - \delta_{3K}}}{2\sqrt{K}(L - 1)} \|\boldsymbol{h}_{T_{0}}\|_{2} \triangleq \beta \|\boldsymbol{h}_{T_{0}}\|_{2}$$
(31)

Since  $\{\boldsymbol{h}_{T_l}\}_{l=1}^{L}$  are arranged in descending order of magnitude, we have

$$|h_i| \ge |h_{T_2^1}| \ge \beta \|\boldsymbol{h}_{T_0}\|_2 \qquad \forall i \in T_1$$
(32)

The proof is completed here.

# APPENDIX II PROOF OF THE INEQUALITY (25)

We relax the terms on both sides of (25) as follows:

$$\prod_{i \in S_2} \left( \frac{|\alpha_i - \eta_i| + 1}{\alpha_i + 1} \right) > \prod_{i \in S_2} \left( \frac{1}{\alpha_i + 1} \right)$$
(33)  
$$\left( \frac{1}{\beta \eta_{\max} + 1} \right)^{2K - k_0} < \left( \frac{1}{\beta \eta_{\max}} \right)^{2K - k_0}$$
$$\stackrel{(a)}{\leq} \frac{1}{\eta_q^K \beta^{2K - k_0}} \prod_{i \in S_2} \frac{1}{\eta_i}$$
$$\stackrel{(b)}{=} \frac{1}{\alpha_q^K \gamma_q^K \beta^{2K - k_0}} \prod_{i \in S_2} \frac{1}{\alpha_i \gamma_i}$$
(34)

where  $\eta_q \triangleq \min_{i \in S_2} \eta_i$ , (a) comes from the fact:  $\eta_{\max} \ge \eta_i$ ,  $\forall i \in T_0$  and the cardinality of  $S_2$  is equal to  $K - k_0$ , and in (b), we replace  $\eta_i$  with  $\alpha_i \gamma_i$  as we have  $\gamma_i = \eta_i / \alpha_i$ . To prove (25), it suffices to show the following inequality holds valid

$$\prod_{i \in S_2} \left( \frac{1}{\alpha_i + 1} \right) \ge \frac{1}{\alpha_q^K \gamma_q^K \beta^{2K - k_0}} \prod_{i \in S_2} \frac{1}{\alpha_i \gamma_i}$$
(35)

Rearranging (35), we get

$$\alpha_q^K > \varphi \triangleq \frac{1}{\gamma_q^K \beta^{2K-k_0}} \prod_{i \in S_2} \frac{\alpha_i + 1}{\alpha_i \gamma_i}$$
(36)

Note that  $\gamma_i > \tau, \forall i \in S_2$  and

$$\frac{\alpha_i + 1}{\alpha_i} \le \frac{\alpha_{\min} + 1}{\alpha_{\min}} \triangleq d \tag{37}$$

where the above inequality comes from the fact that  $\alpha_i > 0$  and  $\alpha_{\min} \le \alpha_i$ . Hence  $\varphi$  defined in (36) is upper bounded by

$$\varphi < \left(\frac{1}{\beta\tau}\right)^{2K-k_0} d^{K-k_0} < \left(\frac{1}{\beta\tau}\right)^{2K} d^K \tag{38}$$

where the last inequality comes by noting that  $0 < \beta < \frac{1}{2}$  (c.f. (4)),  $0 < \tau < 1$ , and d > 1. Therefore the following

$$\alpha_{\min}^{K} > \left(\frac{1}{\beta\tau}\right)^{2K} d^{K} \Leftrightarrow \alpha_{\min} > d\left(\frac{1}{\beta\tau}\right)^{2}$$
(39)

is a sufficient condition for the inequality (36) (consequently (25)) to hold valid. Moreover, note that the condition (39) implies that  $\alpha_{\min} > 1$ . Therefore we have  $d = (\alpha_{\min}+1)/\alpha_{\min} < 2$ , in which case the condition (39) can be further relaxed as

$$\alpha_{\min} > \frac{2}{\left(\beta\tau\right)^2} \tag{40}$$

The proof is completed here.

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